

NO-ARBITRAGE WITH
NON-SEMIMARTINGALES
CONTINUOUS SIMPLE ARBITRAGE CASE

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ABSTRACT

We discuss some recent notions and results connected to simple arbitrage with general continuous pricing models. The discussion is based on the following works (and other works):

-  Bender, C. (2010) Simple Arbitrage. *Preprint*.
-  Bender, C., Sottinen, T., and Valkeila, E. (2011) Fractional processes as models in stochastic finance. *Advanced Mathematical Methods for Finance* (Eds. G. Di Nunno and B. Oksendal), Springer.
-  Bender, C., Sottinen, T. and Valkeila, E. (2008) Pricing by hedging and no-arbitrage beyond semimartingales. *Finance and Stochastics* **12**, 441–468.

OUTLINE

- 1 SIMPLE ARBITRAGE
- 2 CONDITIONAL FULL SUPPORT
- 3 DELAY-SIMPLE ARBITRAGE

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SIMPLE ARBITRAGE

\mathcal{T} -SIMPLE STRATEGY

$S = (S_t)_{t \in [0, T]}$ is a continuous **DISCOUNTED STOCK-PRICE PROCESS** on a filtered probability space $(\Omega, \mathcal{F}, (\mathcal{F}_t)_{t \in [0, T]}, \mathbf{P})$.

DEFINITION (\mathcal{T} -SIMPLE STRATEGY)

A **\mathcal{T} -SIMPLE STRATEGY** is a trading strategy where the number of stocks owned by the investor at time t is

$$\Phi_t = \varphi_0 \mathbf{1}_{\{0\}}(t) + \sum_{j=0}^{n-1} \varphi_j \mathbf{1}_{(\tau_j, \tau_{j+1}]}$$

φ_j are \mathcal{F}_{τ_j} -measurable, $n \in \mathbb{N}$ is fixed, $\tau_0 \leq \tau_1 \leq \dots \leq \tau_n$ are stopping times, satisfying some additional condition \mathcal{T} to be specified later. If there are no additional conditions, we omit the \mathcal{T} .

SIMPLE ARBITRAGE

(0-ADMISSIBLE) \mathcal{T} -SIMPLE ARBITRAGE

The **WEALTH PROCESS** associated with strategy Φ with initial endowment v is

$$V_t(\Phi, v) = v + \sum_{j=0}^{n-1} \Phi_{\tau_{j+1}} (S_{t \wedge \tau_{j+1}} - S_{t \wedge \tau_j}).$$

DEFINITION ((0-ADMISSIBLE) \mathcal{T} -SIMPLE ARBITRAGE)

A \mathcal{T} -Simple strategy is

- a **\mathcal{T} -SIMPLE ARBITRAGE** if $V_T(\Phi, 0) \geq 0$ **P**-a.s. and $\mathbf{P}[V_T(\Phi, 0) > 0] > 0$.
- **c -ADMISSIBLE** if $V_t(\Phi, 0) \geq -c$ for all $t \in [0, T]$ **P**-a.s.

SIMPLE ARBITRAGE

CONDITIONAL UP'N'DOWN (CUD)

DEFINITION (\mathcal{T} -CUD)

The stock-price S satisfies the \mathcal{T} -CONDITIONAL UP'N'DOWN PROPERTY (\mathcal{T} -CUD), if, \mathbf{P} -a.s.,

$$\begin{aligned}\mathbf{P} [S_{\tau_{j+1}} - S_{\tau_j} > 0 | \mathcal{F}_{\tau_j}] &> 0, \text{ and} \\ \mathbf{P} [S_{\tau_{j+1}} - S_{\tau_j} < 0 | \mathcal{F}_{\tau_j}] &> 0,\end{aligned}$$

or $\mathbf{P}[S_{\tau_{j+1}} = S_{\tau_j} | \mathcal{F}_{\tau_j}] = 1$ for all $\tau_j \leq \tau_{j+1}$ satisfying \mathcal{T} .

LEMMA (CUD CHARACTERIZATION OF SIMPLE ARBITRAGE)

The model S admits \mathcal{T} -simple arbitrage if and only if it does not satisfy \mathcal{T} -CUD.

SIMPLE ARBITRAGE

NO OBVIOUS ARBITRAGE (NOA)

DEFINITION (NO OBVIOUS ARBITRAGE (NOA))

S satisfies **NO OBVIOUS ARBITRAGE (NOA)** if for all stopping times τ and $\varepsilon > 0$

$$\mathbf{P} \left[\inf_{t \in [\tau, T]} X_t - X_\tau > -\varepsilon \right] > 0 \text{ and } \mathbf{P} \left[\sup_{t \in [\tau, T]} X_t - X_\tau < \varepsilon \right] > 0$$

REMARK

If NOA is violated, then there exists 0-admissible **OBVIOUS ARBITRAGE** of the type $\Phi = \pm \mathbf{1}_{[\tau_1, \tau_2]}$.

SIMPLE ARBITRAGE

INFINITESIMAL UP'N'DOWN (IUD)

DEFINITION (INFINITESIMAL UP'N'DOWN (IUD))

Let τ be a stopping time. Let

$$\tau^\pm = \inf\{t \geq \tau; \pm S_t > \pm S_\tau\}.$$

S satisfies **INFINITESIMAL UP'N'DOWN (IUD)** if $\tau^+ = \tau^-$ **P**-a.s. for all stopping times.

THEOREM (NO SIMPLE ARBITRAGE, NOA AND IUD)

- 1** S satisfies IUD iff S does not admit 0-admissible simple arbitrage.
- 2** S satisfies IUD and NOA iff S does not admit simple arbitrage.

SIMPLE ARBITRAGE

MODELS WITH NOA AND IUD

Checking if a model satisfies NOA seems very difficult, since it involves stopping times. Ditto for IUD. There is the following though:

THEOREM (BENDER 2010)

Let $S = M + Y$, where M is a continuous martingale with bracket $\langle M \rangle$ and Y is $1/2$ -Hölder continuous w.r.t. $\langle M \rangle$. Then S satisfies IUD.

SKETCH OF PROOF.

By time-change, consider the Brownian case $M = W$, and use the strong Markov property together with the law of iterated logarithm.

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CONDITIONAL FULL SUPPORT

DEFINITION

DEFINITION (CONDITIONAL FULL SUPPORT (CFS))

S satisfies **CONDITIONAL FULL SUPPORT (CFS)** if, \mathbf{P} -a.s.,

$$\text{supp}(\text{Law}((S_u)_{u \in [t, T]} | \mathcal{F}_t)) = C_{S_t}^+([t, T])$$

REMARK

- In CFS one can replace t by a stopping time τ !!!
- There are host of theorems for checking CFS for given models.
- CFS implies NOA.
- **UNFORTUNATELY, CFS DOES NOT SEEM TO IMPLY IUD.**

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DELAY-SIMPLE ARBITRAGE

MOTIVATION AND DEFINITION

Since CFS does not (probably) imply IUD, we need to restrict the class of simple strategies. Here the condition \mathcal{T} comes into play.

DEFINITION (CLASS $\mathcal{T}_{\text{delay}}$ OF STOPPING TIMES)

For any stopping time τ , let $C_{S_\tau}^+([\tau, T])$ be the random space of continuous positive paths $\omega = (\omega_t)_{t \in [\tau(\omega), T]}$ with $\omega_{\tau(\omega)} = S_{\tau(\omega)}(\omega)$ fixed.

A sequence of non-decreasing stopping times $\tau = (\tau_j)_{j=0}^n$ satisfies the **DELAY** property $\mathcal{T}_{\text{delay}}$ if for all τ_k there is an \mathcal{F}_{τ_j} -measurable open *delay set* $U_k \subset C_{S_{\tau_j}}^+([\tau_j, T])$ and an \mathcal{F}_{τ_j} -measurable a.s. positive random variable ε_k such that $\tau_{j+1} - \tau_j \geq \varepsilon_j$ in the set $U_j \cap \{\tau_{j+1} > \tau_j\}$.

DELAY-SIMPLE ARBITRAGE

EXAMPLES OF STOPPING TIMES WITH OR WITHOUT DELAY

EXAMPLE (POSITIVE EXAMPLE)

Let a^j, b^j be continuous functions with $a_{\tau_j}^j < 0 < b_{\tau_j}^j$ and let

$$\tau_{j+1} = \inf \left\{ t > \tau_j; S_t - S_{\tau_j} \leq a_t^j \text{ or } S_t - S_{\tau_j} \geq b_t^j \right\}.$$

EXAMPLE (NEGATIVE EXAMPLE)

Let $(\tau_1, \tau_2) = (0, \tau)$ with

$$\tau = \inf \left\{ t > 0; S_t = e^{W_t + t^a} = 1 \right\}$$

for some $a < 1/2$. By the law of iterated logarithm $\tau > 0$, but any open set in $C_1^+([0, T])$ contains sequences (ω^n) with $\tau(\omega^n) \rightarrow 0$.

DELAY-SIMPLE ARBITRAGE

DELAY AND LOCAL CONTINUITY

DEFINITION (LOCAL CONTINUITY (LC))

Let \mathcal{X} and \mathcal{Y} be metric spaces. A function $f : \mathcal{X} \rightarrow \mathcal{Y}$ is **LOCALLY CONTINUOUS** (LC) if for all $x \in \mathcal{X}$ there exists an open $U_x \subset \mathcal{X}$ such that $x \in \bar{U}_x$ and $\lim f(x_n) = f(x)$ whenever $\lim x_n = x$ in U_x .

f is **LOCALLY LOWER-SEMICONtinuous** (LLSC), if in the above, $\liminf f(x_n) \geq f(x)$.

REMARK

LC at x is continuity from the “direction” U_x . However, LC is not directional continuity in the classical sense. If $x \in U_x$ then LC is classical continuity.

DELAY-SIMPLE ARBITRAGE

DELAY AND LOCAL CONTINUITY, CONT.

EXAMPLE

A functional $\tau : C_{S_0}^+([0, T]) \rightarrow [0, T]$ defined by

$$\tau(\omega) = \min \{t; \omega(t) = c\}$$

is LC. Indeed, for $\omega_0 \in C_{S_0}^+([0, T])$, take

$$U_{\omega_0} = \{\omega; \omega(t) > \omega_0(t) \text{ for all } t \in [0, T]\}.$$

LEMMA (LOCAL LOWER SEMICONTINUITY AND DELAY)

If a stopping time τ is LLSC, then it has the delay property.

DELAY-SIMPLE ARBITRAGE

CFS AND DELAY-SIMPLE ARBITRAGE

THEOREM

CFS implies no delay-simple arbitrage.

PROOF.

We need to show that the $\mathcal{T}_{\text{delay}}$ -CUD is satisfied. We may assume $\tau_{j+1} > \tau_j$. We show that $\mathbf{P}[S_{\tau_{j+1}} > S_{\tau_j} | \mathcal{F}_{\tau_j}] > 0$ a.s.; the proof for $\mathbf{P}[S_{\tau_{j+1}} < S_{\tau_j} | \mathcal{F}_{\tau_j}] > 0$ a.s. follows analogously.

By the CFS it is enough to show that $\{S_{\tau_{j+1}} > S_{\tau_j}\} \subset C_{S_{\tau_j}}^+([\tau_j, T])$ contains an open set. Let U_j be an ε_j -delay set for τ_j .

We first assume that U_j contains a strictly increasing paths ω^0 on $[\tau_j, T]$. Denote by $B_{\omega^0}(\eta_j)$ the open η_j -ball around ω^0 . Choosing η_j sufficiently small we have $B_{\omega^0}(\eta_j) \subset U_j$ and $\omega_{\tau_j+\varepsilon_j}^0 > \omega_{\tau_j}^0 + \eta_j$.

DELAY-SIMPLE ARBITRAGE

CFS AND DELAY-SIMPLE ARBITRAGE, CONT.

PROOF, CONT.

Hence,

$$\begin{aligned}\omega_{\tau_{j+1}}(\omega) - S_{\tau_j} &> \omega_{\tau_{j+1}}^0(\omega) - \eta_j - S_{\tau_j} \\ &\geq \omega_{\tau_j + \varepsilon_j}^0 - S_{\tau_j} - \eta_j \\ &= \omega_{\tau_j + \varepsilon_j}^0 - \omega_{\tau_j}^0 - \eta_j \\ &> 0,\end{aligned}$$

So, $B_{\omega^0}(\eta_j) \subset \{S_{\tau_{j+1}} > S_{\tau_j}\}$, and the claim follows, if U_j contains a strictly increasing paths.

If U_j does not contain a strictly increasing path, we proceed as follows:

DELAY-SIMPLE ARBITRAGE

CFS AND DELAY-SIMPLE ARBITRAGE, CONT., CONT.

PROOF, CONT., CONT.

Being an open set in $C_{S_{\tau_j}}^+([\tau_j, T])$, U_j contains paths that are strictly increasing on a small enough interval $[\tau_j, \tau_j + 2\eta_j]$.

Hence, there is a strictly increasing path ω^0 and an open ball B_j around ω^0 in $C_{S_{\tau_j}}^+([\tau_j, T])$ such that any $\omega \in B_j$ coincides with some path $\bar{\omega} \in U_j$ on the segment $[\tau_j, \tau_j + \eta_j]$.

Hence, $\tau_{j+1}(\omega) - \tau_j \geq (\tau_{j+1}(\bar{\omega}) - \tau_j) \wedge \eta_j \geq \epsilon_j \wedge \eta_j =: \epsilon_j^0$ for every $\omega \in B_j$.

Therefore B_j is an ϵ_j^0 -delay set which contains a strictly increasing path and so the first case applies. \square